











Sommaire

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Managing code variability in HPC applications







Managing code variability in HPC applications High-level code features to be updated Low-level code stencil . FFT . **StarPU** OpenMP CUDA new runtime MPI





Managing code variability in HPC applications High-level code High-level code features to be updated Low-level code Low-level code stencil stencil . FFT FFT . **StarPU** new application OpenMP CUDA OpenMP CUDA new runtime MPI **StarPU** MPI











Component-based programming model







Component-based programming model







Component-based programming model







L2C and COMET

L2C

- Ports : Use + Provide
- Ports: Corba + MPI
- Abstract C++ classes to define services
- Annotated C++ classes to define components
- XML assembly description file





L2C and COMET







COMET composition units







COMET composition units



Metatask description





COMET composition units







COMET assembly and a possible generated task graph







Implementing the COMET model











CIGRA : COMET Instruction GRAph







Reverse topological sorting to a CIGRA

























Kinds of dynamic data connection with COMET







Flat code generation strategy







Flat code generation strategy







Nested code generation strategy



#pragma omp parallel #pragma omp single #pragma omp task o(A) A.submit(); #pragma omp taskwait Update C size #pragma omp task io(C) i(A) C.submit(); #pragma omp taskwait #pragma omp task io(E) i(C) E.submit(): #pragma omp taskwait #pragma omp task o(B) B.submit(); #pragma omp taskwait Update B size





Weak-nested code generation strategy



#pragma omp parallel
{
 #pragma omp single
 {
 #pragma omp task o(A) weak
 A.submit();
 #pragma omp taskwait
 Update C size
 #pragma omp task io(C) i(A) weak
 C.submit();
 #pragma omp task io(E) i(C) weak
 E.submit();
 #pragma omp task o(B) weak
 B.submit();
 #pragma omp taskwait
 Update B size
 }
...























TR1

EP2

Gen2

EP1

TR2

Gen1

2. Handling Dynamic Data-Parallel Dataflows with COMET

Gen1

EP1

EP3

EP2

EP4

Gen2

Experimental results on Grid'5000 on the roazhon4 node Datasize : 8192x8192

Blocksizes :

1024x1024, 1024x512, 1024x256, 1024x128 **Runtimes :**

llvm-openmp, mpc-openmp, llvm-ompss2







Experimental results on Grid'5000 on the roazhon4 node

Observations :

- In microbenchs, flat is better than nested and weak-nested whatever the runtime.
- In mB1 and mB2, weak-nested is better than flat and nested on the llvm-ompss2 runtime.
- In mB1 and mB2, nested is than flat and nested on llvm-ompss2 runtime.
- There are some anomalies...













3. Conclusion

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3. Conclusion

Wrap-up :

- COMET is a component-based programming model designed for creating HPC applications through fine-grained, dataflow-like compositions.
- We have extended it to manage applications designed by dynamic data-parallel dataflows.
- Experiments on several benchmarks show that the execution times of our different code generation approaches depend on the runtime used.
- OpenMP needs to be upgraded to incorporate the concept of weak dependencies.

What's next ?

- Understand and correct any anomalies observed.
- We have done a micro-aevol version with COMET but the performance is not good. We need to do something about that.
- We have started integrating StarPU to COMET. We are working on the management of partitioned data between both.







